

# ALTERNATIVE SOLUTIONS TO DILUTED P-SPIN MODELS AND RANDOM XORSAT PROBLEMS

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  - Solution with rigorous methods
  - Generalized XORSAT
- } comparison
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Independent work on rigorous results by  
Cocco, Dubois, Mandler and Monasson

## Motivations

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Original motivation: simplify random 3-SAT keeping relevant features of the phase transition.

$$\mathcal{H}_{3\text{SAT}} \propto \alpha N - \sum_{i=1}^N H_i s_i + \sum_{ij} T_{ij} s_i s_j - \sum_{ijk} J_{ijk} s_i s_j s_k$$

↓

$$\mathcal{H}_{3\text{XORSAT}} \propto \alpha N - \sum_{ijk} J_{ijk} s_i s_j s_k$$

- Exactly solvable model with a hard-SAT phase, very similar to random K-SAT.
- Show correctness of replica and cavity calculations on a non-trivial model.

## Definition of the model

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1) Diluted  $p$ -spin model at  $T = 0$

$$\mathcal{H} = \sum_{\{i,j,k\} \in G} (1 - J_{ijk} s_i s_j s_k)$$

$G$ : set of  $\alpha N$  triples randomly chosen among  $\binom{N}{3}$   
(random hypergraph with average degree  $3\alpha$ )

$J_{ijk} = \pm 1$  quenched random variables

**Find  $s_i = \pm 1$  such that  $s_i s_j s_k = J_{ijk} \quad \forall \{i,j,k\} \in G$**

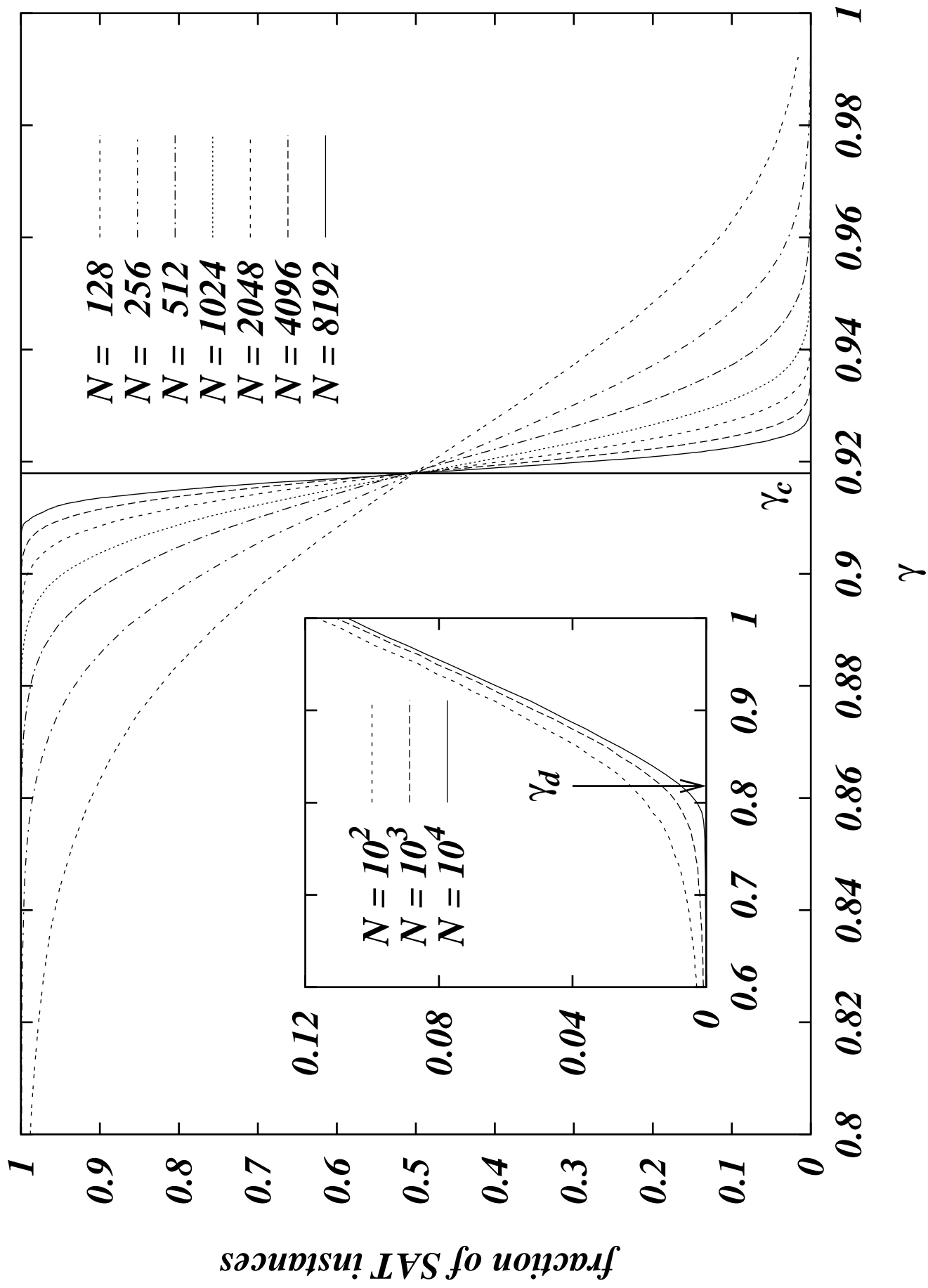
2) Set of  $\alpha N$  linear equations mod 2 in  $N$  variables

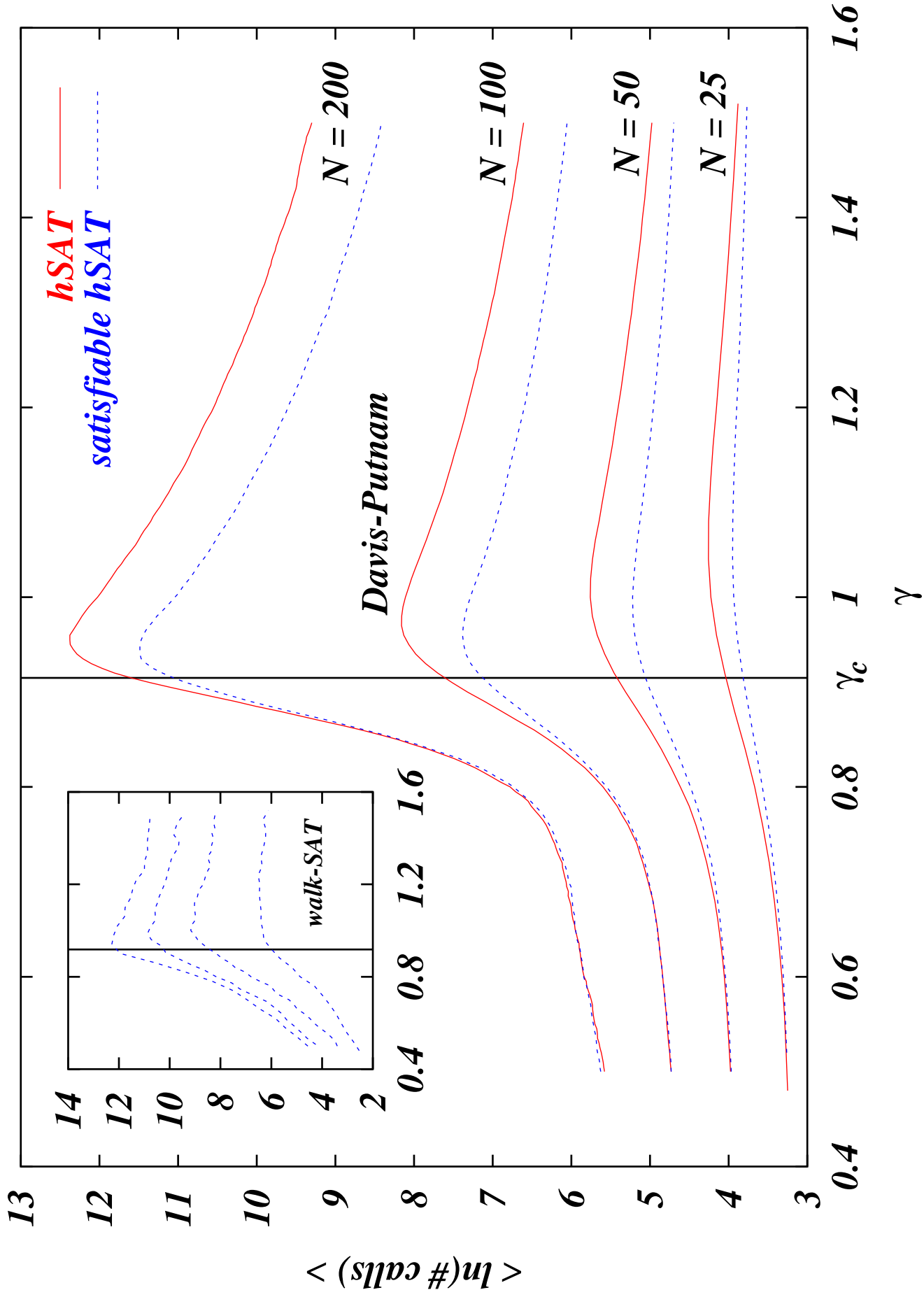
$$s_i = (-1)^{x_i} \quad J_{ijk} = (-1)^{y_{ijk}} \quad x, y \in \{0, 1\}$$

$$s_i s_j s_k = J_{ijk} \quad \iff \quad x_i + x_j + x_k = y_{ijk} \pmod{2}$$

3) Random XORSAT problem:

$$F = \bigwedge_{\{i,j,k\} \in G} x_i \oplus x_j \oplus x_k \oplus y_{ijk}$$





## Questions

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Can we locate the phase transition?

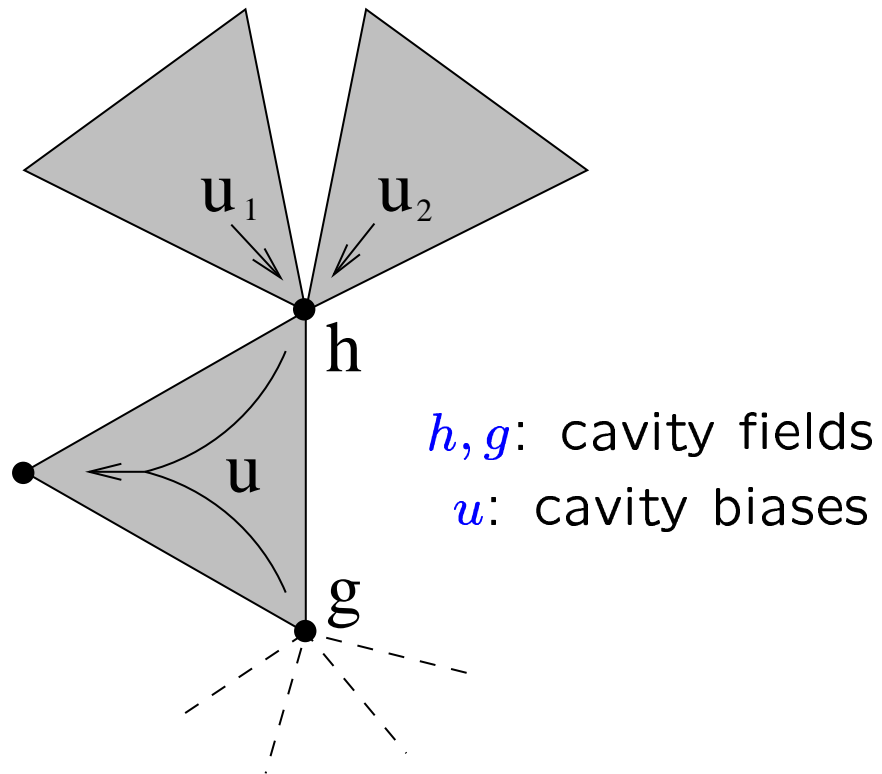
Which kind of phase transition is it?

Is there any hard-SAT region?

Why finding solutions is hard in this region?

## Solution with cavity

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$$P(h) \propto \int \prod_{i=1}^k dQ_i(u_i) \delta(h - \sum u_i) e^{-\mu(\sum |u_i| - |\sum u_i|)}$$

$$Q(u) = \int dP(h) dP(g) \delta(u - \text{sign}(hg))$$

## Solution with cavity

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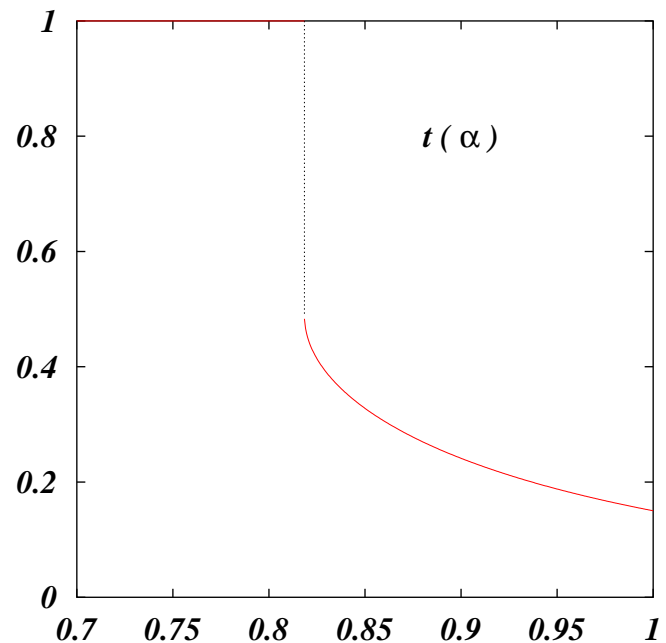
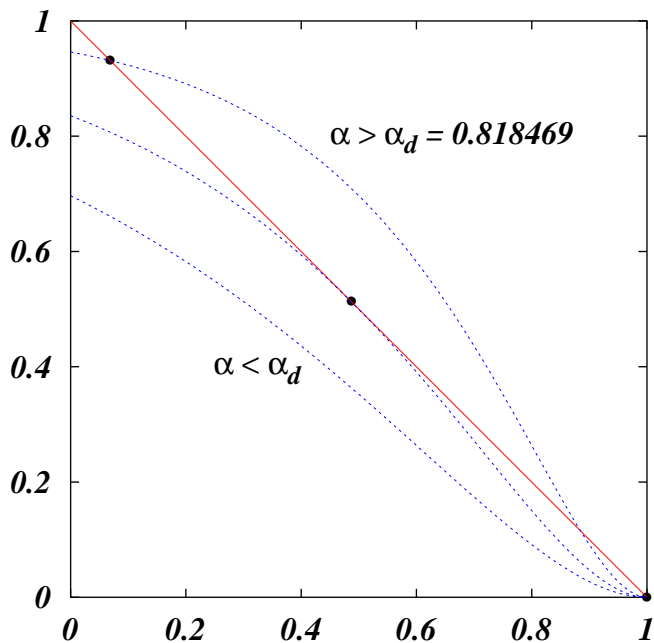
For  $\mu \rightarrow \infty$

$$Q(u) = \begin{cases} \delta(u) & \text{prob. } t \\ \frac{1}{2}[\delta(u+1) + \delta(u-1)] & \text{prob. } 1-t \end{cases}$$

Prob.  $P(h = \sum_{i=1}^k u_i)$  is non trivial =  $1 - t^k$

$$e^{-3\alpha} \sum_{k=0}^{\infty} \frac{(3\alpha)^k}{k!} (1 - t^k) = 1 - e^{-3\alpha(1-t)}$$

$$1 - t = \left(1 - e^{-3\alpha(1-t)}\right)^2$$





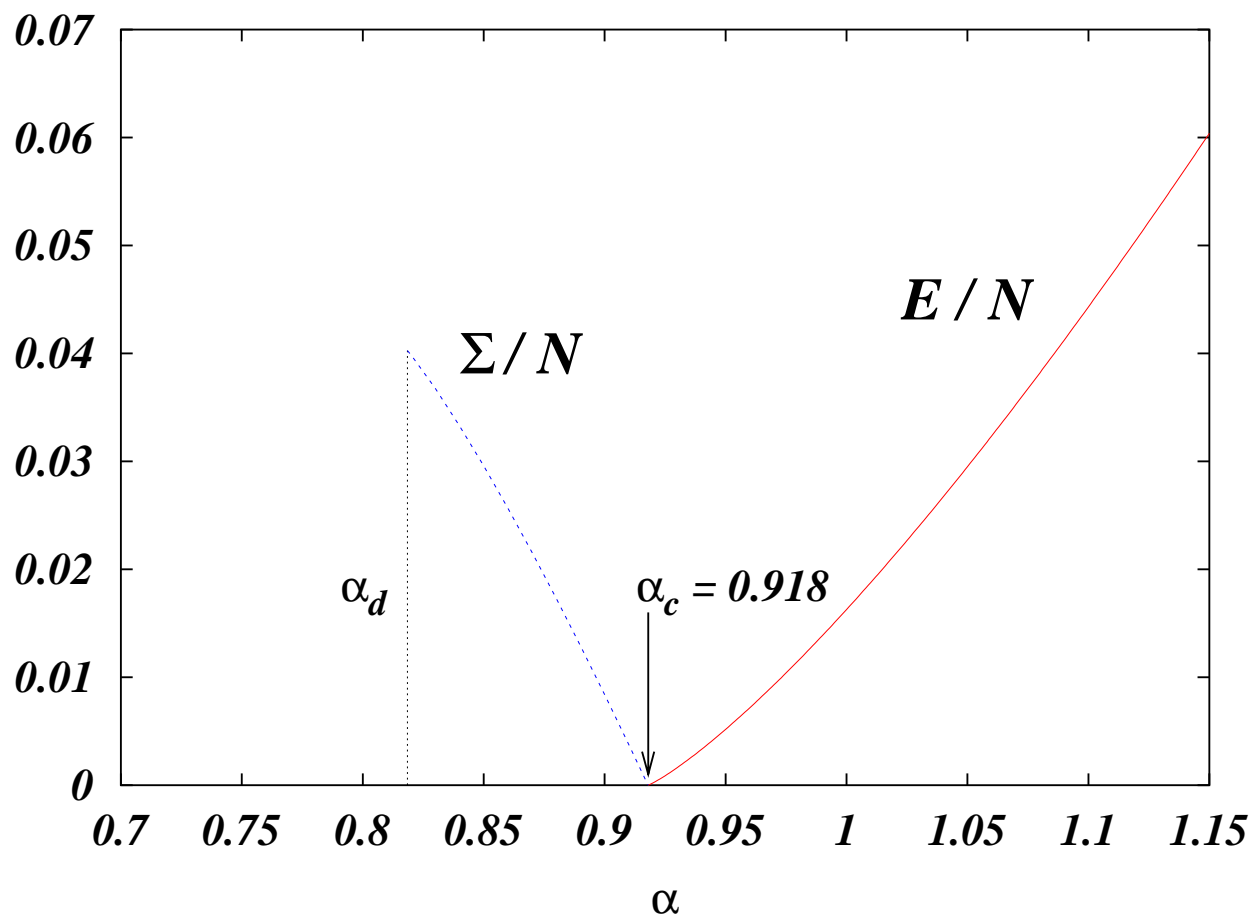
# Solution with cavity

$$\lambda = 3\alpha(1 - t) \quad \lambda = 3\alpha(1 - e^{-\lambda})^2$$

For  $\mu \gg 1$   $\Phi(\mu) \simeq \frac{1}{\mu} (\psi - \omega e^{-2\mu})$  with

$$\psi = \log(2) \left[ \frac{\lambda}{3} - 1 + e^{-\lambda} \left( 1 + \frac{2}{3}\lambda \right) \right]$$

$$\omega = \frac{\lambda}{6} [2 - e^{-\lambda}(2 + 3\lambda)]$$



# Rigorous solution

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## Physical idea

The hypergraph can be divided in 2 parts:

- 1) a central core, where  $Q(u)$  are non-trivial and variables are more constrained;
- 2) an external part, made of dangling ends, where  $Q(u)$  are trivial and variables much less constrained.

Variables in the non-core part induce large fluctuations in the number of solutions  $\mathcal{N}$ , such that

$$\overline{\log(\mathcal{N})} \neq \log(\overline{\mathcal{N}})$$

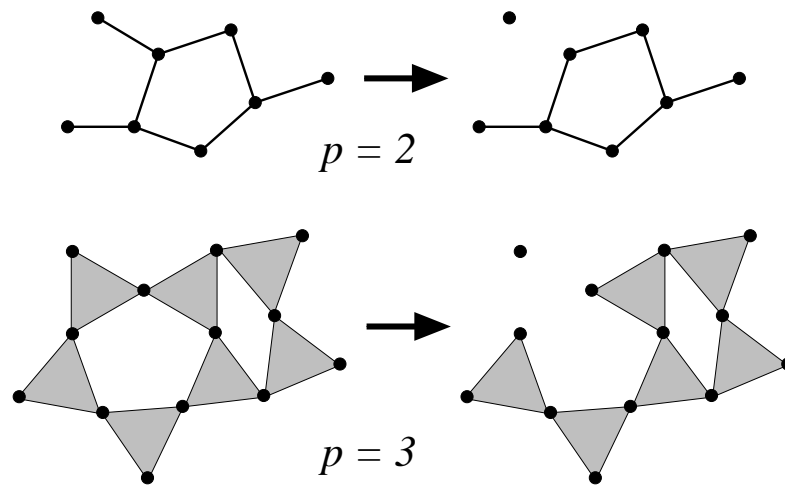
## Plan

1. remove these variables (low connectivity)
2. calculate annealed averages  $\overline{\mathcal{N}}$  and  $\overline{\mathcal{N}^2}$  on the remaining hypergraph

## Leaf removal algorithm

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Rule: As long as there are variables of connectivity less than 2, remove them.



$$\frac{\partial f_0(t)}{\partial t} = (p-1) \frac{f_1(t)}{m(t)} + 1$$

$$\frac{\partial f_1(t)}{\partial t} = (p-1) \frac{2f_2(t) - f_1(t)}{m(t)} - 1$$

$$\frac{\partial f_k(t)}{\partial t} = (p-1) \frac{(k+1)f_{k+1}(t) - kf_k(t)}{m(t)}$$

where  $m(t) = \sum_k k f_k(t) = p(\alpha - t)$ .

$$f_k(t) = e^{-\lambda(t)} \frac{\lambda(t)^k}{k!} \quad \forall k \geq 2$$

## Leaf removal algorithm

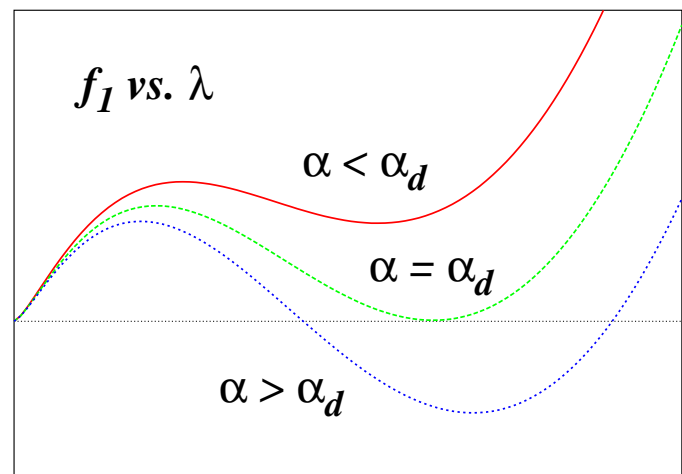
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Solution for  $p = 3$

$$\lambda(t) = 3 \sqrt[3]{\alpha(\alpha - t)^2} \quad f_1(\lambda) = \lambda \left( e^{-\lambda} - 1 + \sqrt{\frac{\lambda}{3\alpha}} \right)$$

First order  
transition at  
 $\alpha = \alpha_d = 0.818469$

The core appears in  
a discontinuous way



On the core:

$$f_0 = f_1 = 0 \quad f_k = \frac{\lambda^k}{k!} e^{-\lambda} \quad \forall k \geq 2$$

$$\text{where } \lambda \text{ solves } \lambda = 3\alpha (1 - e^{-\lambda})^2$$

# On the core

$\mathcal{N}_c$ : number of solution on the core

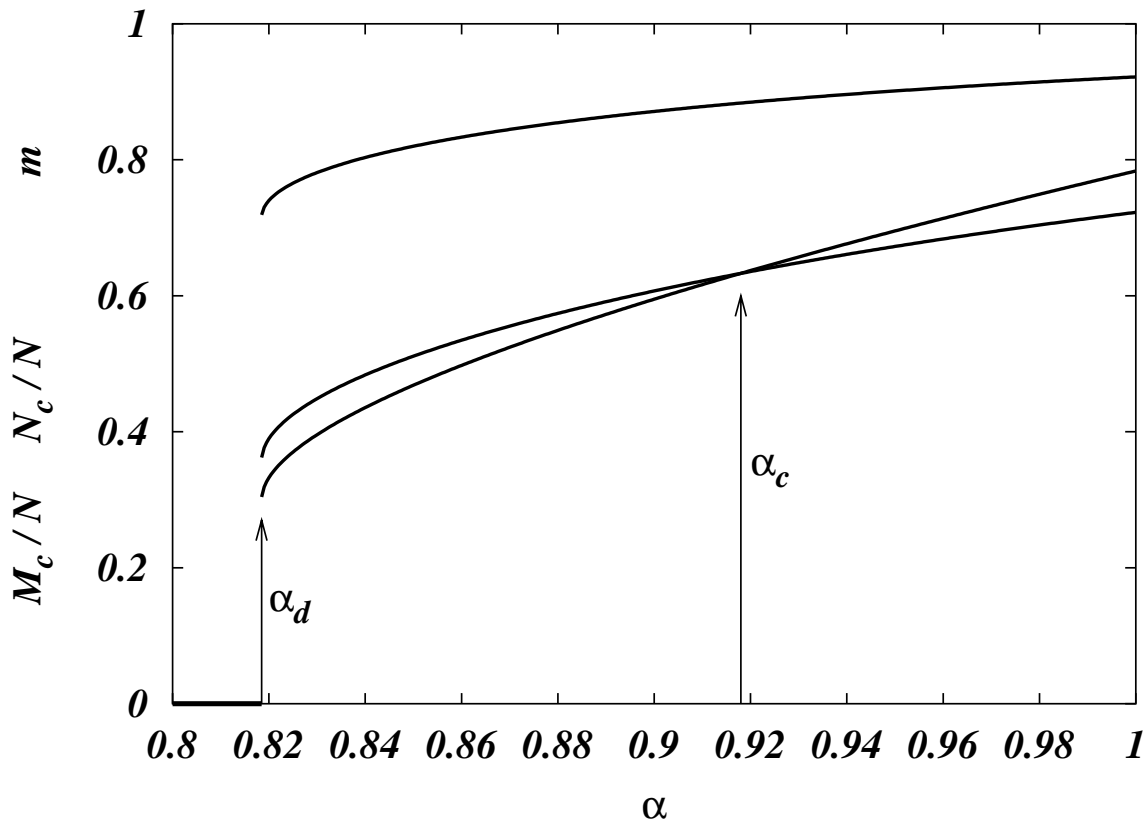
For  $N \rightarrow \infty$ ,

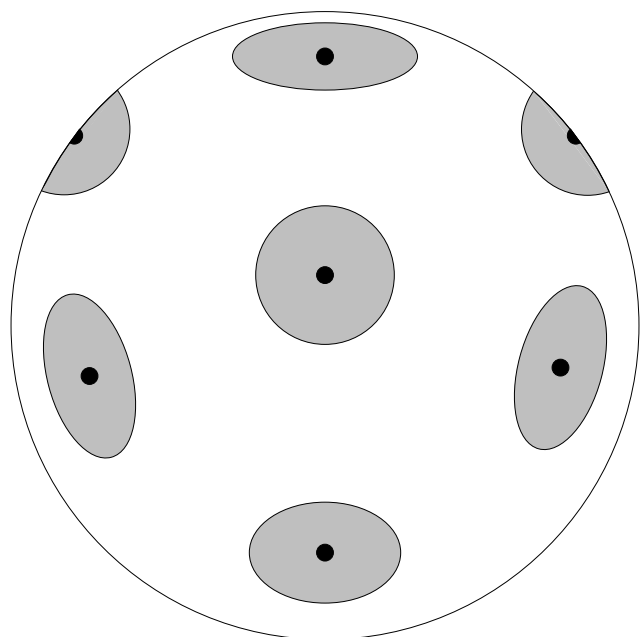
$$\overline{\mathcal{N}_c} = 2^{N_c - M_c} \quad , \quad \frac{\overline{\mathcal{N}_c^2} - \overline{\mathcal{N}_c}^2}{\overline{\mathcal{N}_c}^2} \rightarrow 0$$

$$\overline{\log \mathcal{N}_c} = \log \overline{\mathcal{N}_c} = \log(2)(N_c - M_c)$$

$$N_c(\alpha) = N \left[ 1 - (1 + \lambda)e^{-\lambda} \right]$$

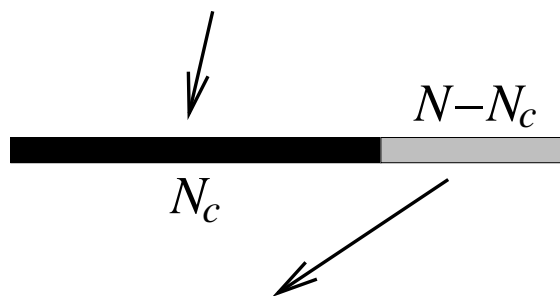
$$M_c(\alpha) = N \frac{\lambda}{p} (1 - e^{-\lambda})$$





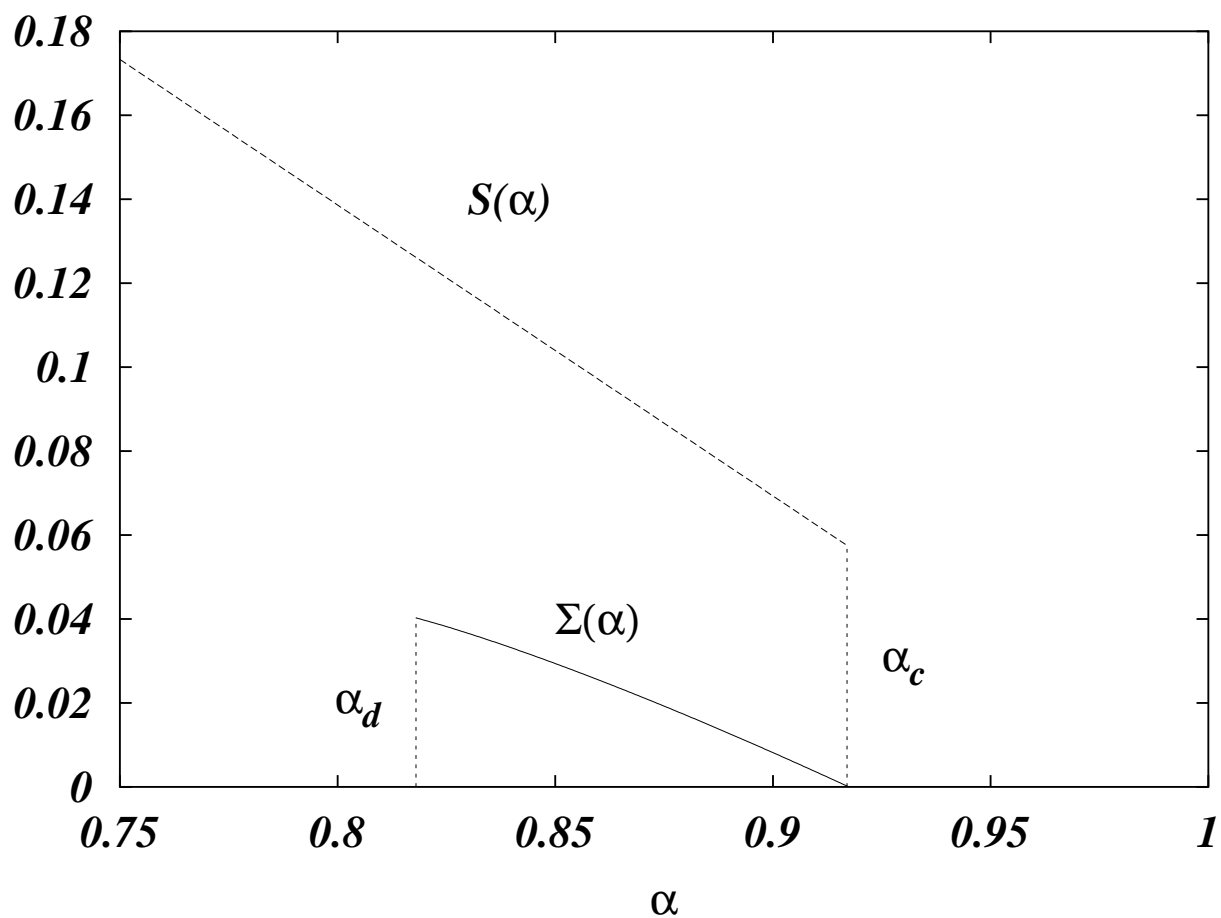
$$\# \text{ clusters} = e^{N\Sigma} = \# \bullet = 2^{N_c - M_c}$$

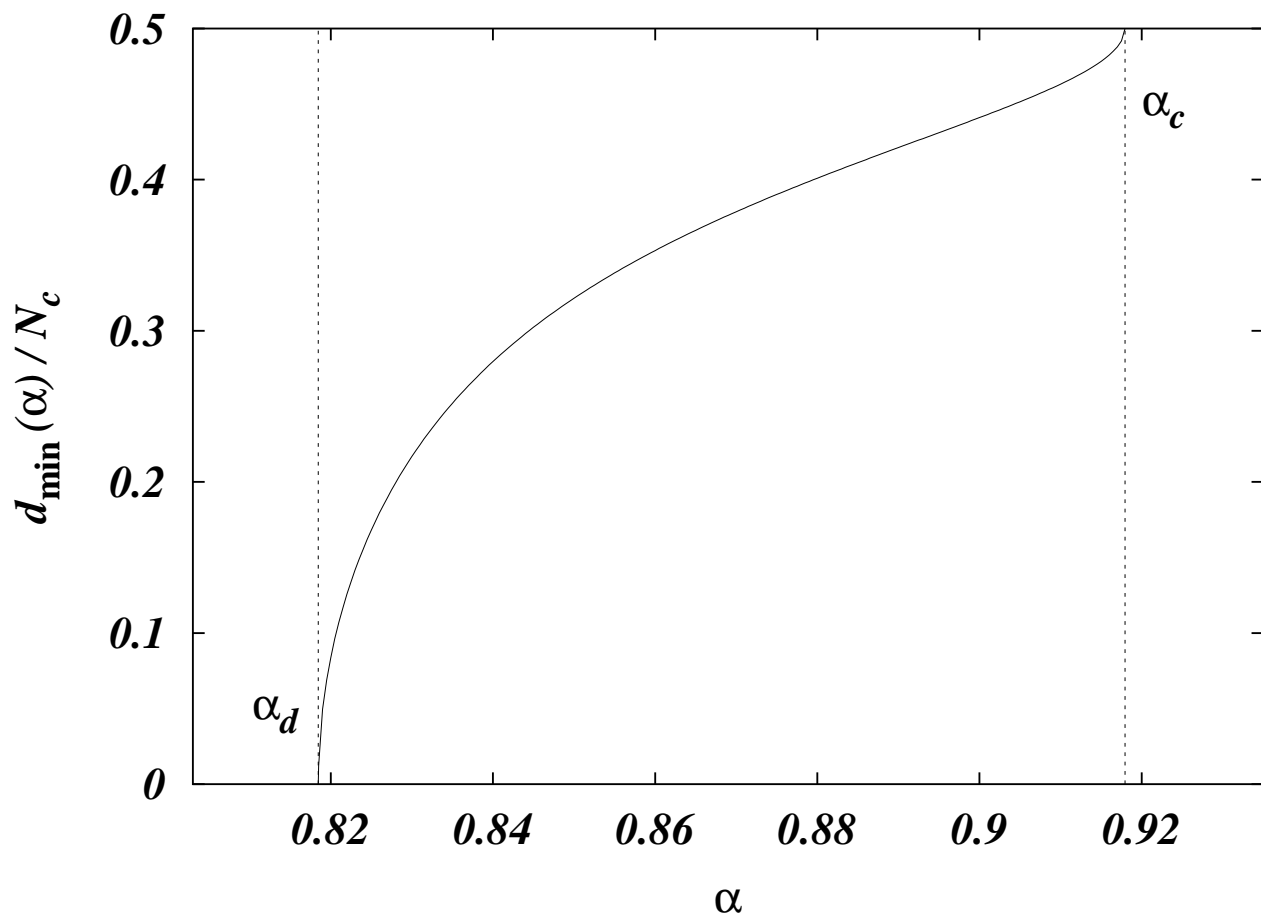
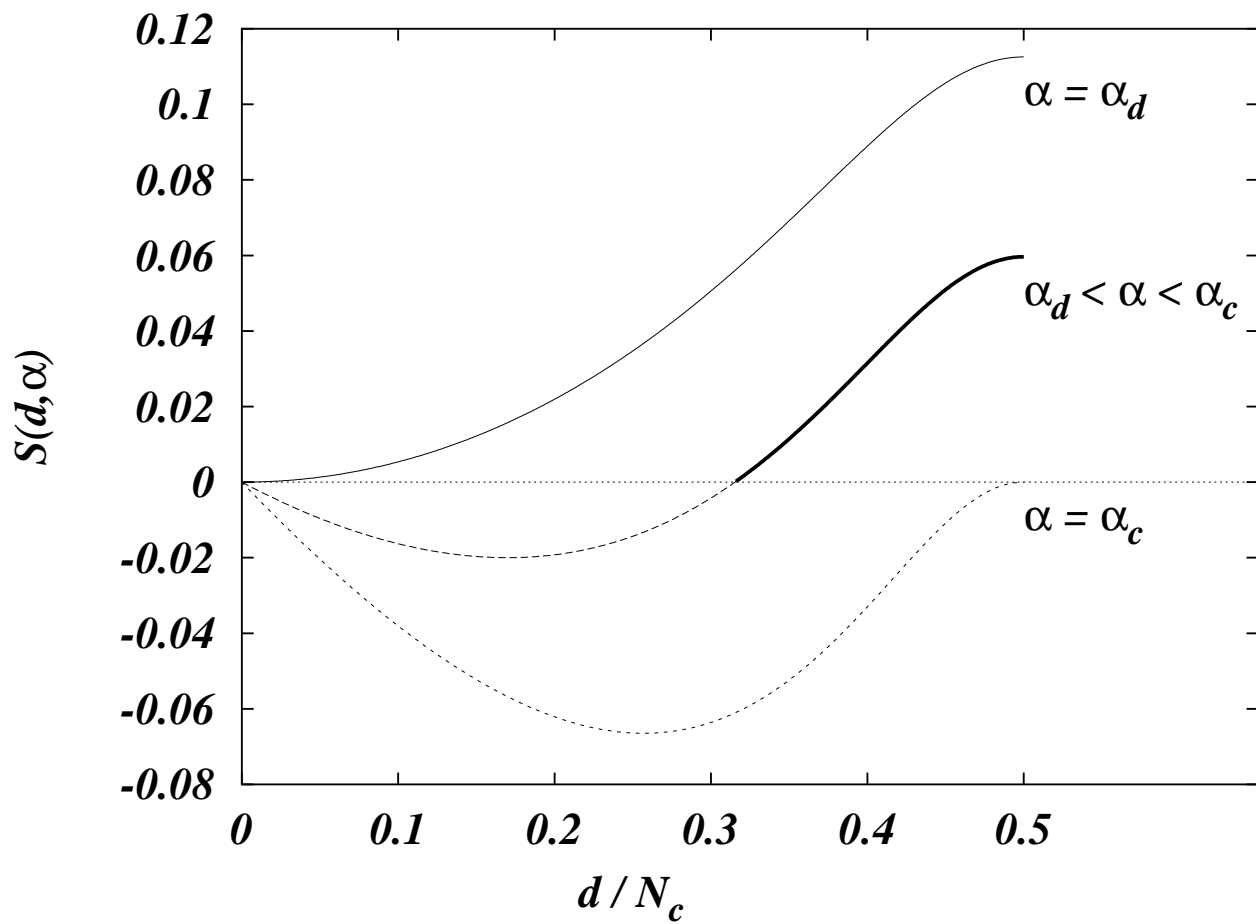
• = core solution



non-core variables give the intra-cluster entropy

$$S_{nc} = S - S_c = S - \Sigma$$





## Generalized XORSAT

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$f_k$ : fraction of variables with conn.  $k$

$g_p$ : fraction of clauses with  $p$  variables ( $g_0 = g_1 = 0$ )

$$f(z) = \sum_k f_k z^k \quad g(z) = \sum_p g_p z^p \quad f'(1) = \alpha g'(1)$$

$\alpha_d$  is determined by the appearance of positive solutions to the equation

$$1 - x = \frac{f' \left[ 1 - \frac{g'(x)}{g'(1)} \right]}{f'(1)} \quad (1)$$

$\alpha_c$  is determined by the equation

$$\alpha g(x^*) = 1 - \alpha (1 - x^*) g'(x^*) - f \left[ 1 - \frac{g'(x)}{g'(1)} \right]$$

where  $x^*(\alpha)$  is the largest solution to Eq.(1)